

# Stochastic Processes: Lecture 5

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# Determinantal point processes

- ▶ A determinantal point process (DPP) on  $\mathbb{R}^D$  is determined by a kernel  $K(x, x')$ .
- ▶ The joint intensities can be written

$$\det \begin{pmatrix} K(x_i, x_i) & K(x_i, x_j) \\ K(x_i, x_j) & K(x_j, x_j) \end{pmatrix}$$

- ▶ The kernel defines an integral operator  $\mathcal{K}$  acting on  $L^2(\mathbb{R}^D)$  that is self-adjoint, positive semidefinite and trace class.

# Joint intensities of a DPP

## Definition 1

The joint intensities of a point process  $N$  are functions (if any exist)  $\rho_k : (\mathbb{R}^D)^k \rightarrow [0, \infty)$  for  $k \geq 1$ , such that for any family of disjoint sets  $D_1, \dots, D_k \subset \mathbb{R}^D$ ,

$$E \left( \prod_{i=1}^k N(D_i) \right) = \int_{\prod D_i} \rho_k(x_1, \dots, x_k) dx_1 \dots dx_k.$$

## Definition 2

A point process  $N$  on  $\mathbb{R}^D$  is said to be a DPP with kernel  $K$  if its joint intensities satisfy

$$\rho_k(x_1, \dots, x_k) = \det (K(x_i, x_j))_{1 \leq i, j \leq k}$$

for every  $k \geq 1$  and  $x_1, \dots, x_k \in \mathbb{R}^D$ .

## Permanental point processes

Leibniz' formula for the determinant of a  $k \times k$  matrix  $M$  is

$$\det(M) = \sum_{\sigma \in S_k} \left( \operatorname{sgn}(\sigma) \prod_{i=1}^k M_{i,\sigma(i)} \right).$$

We denote the *permanent* of a  $k \times k$  matrix  $M$

$$\operatorname{per}(M) = \sum_{\sigma \in S_k} \prod_{i=1}^k M_{i,\sigma(i)}.$$

### Definition 3

A point process  $N$  on  $\mathbb{R}^D$  is said to be a *permanental point process* with kernel  $K$  if its joint intensities satisfy

$$\rho_k(x_1, \dots, x_k) = \operatorname{per}(K(x_i, x_j))_{1 \leq i, j \leq k}$$

for every  $k \geq 1$  and  $x_1, \dots, x_k \in \mathbb{R}^D$ .

# Poisson processes, DPPs and PPPs

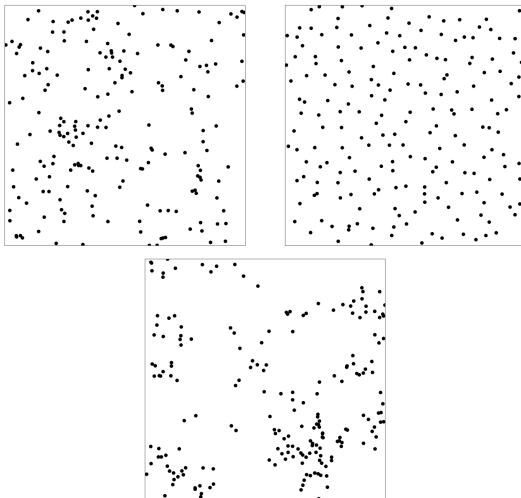


FIG 1. *Samples of translation invariant point processes in the plane: Poisson (left), determinantal (center) and permanental for  $K(z, w) = \frac{1}{\pi} e^{z\bar{w} - \frac{1}{2}(|z|^2 + |w|^2)}$ . Determinantal processes exhibit repulsion, while permanental processes exhibit clumping.*

# DPP results

## Lemma 1

*Suppose  $\{\phi_k\}_{k=1}^n$  is an orthonormal set in  $L^2(\mathbb{R}^D)$ . Then there exists a DPP with kernel*

$$K(x, y) = \sum_{k=1}^n \phi_k(x) \overline{\phi_k(y)}.$$

## Theorem 1

*Let  $K$  determine a self-adjoint integral operator  $\mathcal{K}$  on  $L^2(\mathbb{R}^D)$  that is locally trace-class. Then  $K$  defines a DPP on  $\mathbb{R}^D$  iff all the eigenvalues of  $\mathcal{K}$  are in  $[0, 1]$ .*

# DPP results

## Theorem 2

Suppose  $N$  is a DPP with kernel  $K(x, y)$ . Write

$$K(x, y) = \sum_{k=1}^{\infty} \lambda_k \phi_k(x) \bar{\phi}_k(y),$$

where  $\phi_k$  are normalized eigenfunctions with eigenvalues  $\lambda_k \in [0, 1]$ . Let  $I_k \stackrel{\perp}{\sim} \text{Bernoulli}(\lambda_k)$  and define  $K$ 's random analogue

$$K_I(x, y) = \sum_{k=1}^{\infty} I_k \phi_k(x) \bar{\phi}_k(y).$$

Let  $N_I$  be a DPP with kernel  $K_I$ . Then

$$N \stackrel{d}{=} N_I.$$

In particular, the total number of points in  $N$  follows the distribution of the sum of independent  $\text{Bernoulli}(\lambda_k)$  r.v.s.

## DPP example: non-intersecting random walks

Consider  $n$  independent simple symmetric walks on  $\mathbb{Z}$  started from  $i_1 < \dots < i_n$ , all even. Let  $P_{ij}(t)$  be the  $t$ -step transition probabilities. The probability the r.w.s are at  $j_1 < \dots < j_n$  at time  $t$  and have non-intersecting paths is

$$\det \begin{pmatrix} P_{i_1 j_1}(t) & \dots & P_{i_1 j_n}(t) \\ \vdots & \ddots & \\ P_{i_n j_1}(t) & & P_{i_n j_n}(t) \end{pmatrix}.$$

If  $t$  is even and we condition the walks to return to  $i_1, \dots, i_n$  at time  $t$ , then the positions at time  $t/2$  follow a DPP with Hermitian kernel.



## DPP example: Ginibre ensemble

Let  $Q$  be an  $n \times n$  matrix with i.i.d. complex standard normal entries. The eigenvalues of  $Q$  form a DPP on  $\mathbb{C}$  with the kernel

$$K_n(z, w) = \frac{1}{\pi} e^{-\frac{1}{2}(|z|^2 + |w|^2)} \sum_{k=0}^{n-1} \frac{(z\bar{w})^k}{k!}.$$

As  $n \rightarrow \infty$ , we have a DPP on  $\mathbb{C}$  with kernel

$$\begin{aligned} K(z, w) &= \frac{1}{\pi} e^{-\frac{1}{2}(|z|^2 + |w|^2)} \sum_{k=0}^{\infty} \frac{(z\bar{w})^k}{k!} \\ &= \frac{1}{\pi} e^{-\frac{1}{2}(|z|^2 + |w|^2) + z\bar{w}}. \end{aligned}$$

## Zero set of a Gaussian analytic function

The power series  $f(z) = \sum_{n=0}^{\infty} a_n z^n$ , where  $a_n$  are i.i.d. standard complex normals defines a random analytic function on the unit disk (a.s.). The zero set of  $f$  is a determinantal process in the disk with the Bergman kernel

$$K(z, w) = \frac{1}{\pi(1 - z\bar{w})^2} = \frac{1}{\pi} \sum_{k=0}^{\infty} (k+1)(z\bar{w})^k.$$

## DPPs on discrete sets

Let  $\mathcal{Y}$  be a discrete set with  $n$  items. A point process  $N$  on  $\mathcal{Y}$  is a probability distribution on the power set  $2^{\mathcal{Y}}$ .

### Definition 4

*A point process  $N$  is a determinantal point process if for  $Y \subseteq \mathcal{Y}$  randomly sampled according to  $N$  we have for every  $S \subseteq \mathcal{Y}$*

$$\Pr(S \subseteq Y) = \det K_S$$

*for some similarity matrix  $K \in \mathbb{R}^{n \times n}$  that is symmetric and positive semidefinite.*

Let  $S$  be a two-element set with elements  $i$  and  $j$ . Then

$$\Pr(S \subset Y) = K_{ii}K_{jj} - K_{ij}^2 = \Pr(i \subset Y)\Pr(j \subset Y) - K_{ij}^2.$$

# Conditioning

DPPs are closed under conditioning:

$$\begin{aligned}\Pr(A \subseteq Y | B \subseteq Y) &= \Pr(A \cup B \subseteq Y) / \Pr(A \subseteq Y) \\ &= \frac{\det K_{A \cup B}}{\det K_A} \\ &= \frac{\det(K_A) \det(K_B - K_{BA} K_A^{-1} K_{AB})}{\det(K_A)} \\ &= \det(K_B - K_{BA} K_A^{-1} K_{AB}) \\ &= \det([K - K_{Y^c A} K_A^{-1} K_{A Y^c}]_B) .\end{aligned}$$

## Restrictions on $K$

- ▶ Because marginal probabilities of any set  $S \subseteq \mathcal{Y}$  must be in  $[0, 1]$ , all  $\det(K_S) \geq 0$  and hence  $K$  must be positive semidefinite.
- ▶ Moreover, all eigenvalues of  $K$  must inhabit  $[0, 1]$ , i.e.  $0 \preceq K \preceq 1$ .
- ▶ Any  $K$  satisfying  $0 \preceq K \preceq 1$  defines a DPP.

# L-ensembles

- ▶ L-ensembles provide a convenient way to avoid dealing with  $K \preceq 1$  constraints.
- ▶ An L-ensemble is defined using a symmetric matrix  $L \succeq 0$  that defines the *atomic* probability of an event set  $S$  thus:

$$\Pr_L(S) = \Pr(S = Y) \propto \det(L_Y)$$

- ▶ Conveniently, the normalizing constant is known:

$$\sum_{S \subseteq \mathcal{Y}} \det(L_S) = \det(L + I).$$

# L-ensembles

Theorem 3

For any  $S \subseteq \mathcal{Y}$

$$\sum_{S \subseteq Y \subseteq \mathcal{Y}} \det(L_Y) = \det(L + I_{S^c})$$

Corollary 1

$$\sum_{Y \subseteq \mathcal{Y}} \det(L_Y) = \det(L + I)$$

Proof.

Let  $S$  from Theorem 3 equal the empty set. □

# L-ensembles

## Theorem 4

An L-ensemble is a DPP and its marginal kernel is

$$K = L(L + I)^{-1} = I - (L + I)^{-1}$$

## Proof.

The marginal probability of a set  $S$  under the L-ensemble is

$$\begin{aligned}\Pr_L(S \subseteq Y) &= \frac{\sum_{S \subseteq Y \subseteq \mathcal{Y}} \det(L_Y)}{\sum_{Y \subseteq \mathcal{Y}} \det(L_Y)} = \frac{\det(L + I_{S^c})}{\det(L + I)} \\ &= \det\left((L + I_{S^c})(L + I)^{-1}\right) \\ &= \det\left(I_{S^c}(L + I)^{-1} + I - (L + I)^{-1}\right) \\ &= \det\left(I_{S^c}(L + I)^{-1} + (I_S + I_{S^c})\left(I - (L + I)^{-1}\right)\right) \\ &= \det(I_{S^c} + I_S K) = \begin{vmatrix} I_{|S^c| \times |S^c|} & 0 \\ K_{S, S^c} & K_S \end{vmatrix} = \det(I_{|S^c| \times |S^c|}) \det(K_S) \\ &= \det(K_S).\end{aligned}$$

□



# L-ensembles

- ▶ Given a marginal kernel, we may construct an L-ensemble by setting  $L = K(I - K)^{-1}$ .
- ▶ The inverse of  $I - K$  might not exist, so DPPs are a larger class than L-ensembles.
- ▶ If  $L = \sum_k \lambda_k v_k v_k^T$ , then  $K = \sum_k \frac{\lambda_k}{1 + \lambda_k} v_k v_k^T$ .
- ▶ Linear kernel. Let  $X$  be an  $n \times p$  design matrix (set of feature vectors). Taking  $L = XX^T$ , we have

$$\Pr_L(S) \propto \det(L_S) = \text{Vol}^2(\{x_i\}_{i \in S})$$

If  $p < n$ , the DPP will only have  $p$  points.

## Working with DPPs

- ▶ Complements: if  $Y \sim DPP(K)$ , then  $Y^c \sim DPP(I - K)$
- ▶ Conditioning:

$$\Pr_L(Y = S_{in} \cup B | S_{in} \subseteq Y, S_{out} \cap Y = \emptyset) = \frac{\det(L_{S_{in} \cup B})}{\det(L_{S_{out}^c} + I_{S_{in}^c})}$$

- ▶ Marginalization:

$$\Pr(B \subseteq Y | S \subseteq Y) = \det\left(\left[I - \left[(L + I_{S^c})^{-1}\right]_{S^c}\right]_B\right)$$

- ▶ Scaling: if  $K' = \gamma K$  for  $\gamma \in [0, 1]$ , then for all  $S \subseteq \mathcal{Y}$

$$\Pr_{K'}(S \subseteq Y) = \det(K'_S) = \gamma^{|S|} K_S.$$

# Elementary DPPs

- ▶ A DPP is elementary if every eigenvalue of  $K$  is 0 or 1.
- ▶  $N^V$  denotes an elementary DPP with marginal kernel  $K^V = \sum_{v \in V} vv^T$  if  $V$  is a set of orthonormal vectors.
- ▶ The expected total count for a DPP is

$$E(|Y|) = E\left(\sum_{i=1}^n \mathbf{1}\{i \in Y\}\right) = \sum_{i=1}^n \Pr(i \in Y) = \sum_{i=1}^n K_{ii} = \text{tr}(K).$$

- ▶ For an elementary DPP this is

$$E(|Y|) = \text{tr}(K^V) = \text{tr}\left(\sum_{v \in V} vv^T\right) = \sum_{v \in V} v^T v = |V|.$$

- ▶ Furthermore,  $|Y| = |V|$  a.s. because  $\det(K_Y^V) = 0$  when  $|Y| > |V|$ .

# DPPs as mixtures of elementary DPPs

## Lemma 2

A DPP with kernel  $L = \sum_{i=1}^n \lambda_i v_i v_i^T$  is a mixture of elementary DPPs:

$$Pr_L = \frac{1}{\det(L + I)} \sum_{J \subseteq \{1, 2, \dots, n\}} Pr^{V_J} \prod_{i \in J} \lambda_i$$

where  $V_J = \{v_i\}_{i \in J}$

# Sampling DPPs

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**Algorithm 1** Sampling from a DPP

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**Input:** eigendecomposition  $\{(\mathbf{v}_n, \lambda_n)\}_{n=1}^N$  of  $L$

$J \leftarrow \emptyset$

**for**  $n = 1, 2, \dots, N$  **do**

$J \leftarrow J \cup \{n\}$  with prob.  $\frac{\lambda_n}{\lambda_n + 1}$

**end for**

$V \leftarrow \{\mathbf{v}_n\}_{n \in J}$

$Y \leftarrow \emptyset$

**while**  $|V| > 0$  **do**

    Select  $i$  from  $\mathcal{Y}$  with  $\Pr(i) = \frac{1}{|V|} \sum_{\mathbf{v} \in V} (\mathbf{v}^\top \mathbf{e}_i)^2$

$Y \leftarrow Y \cup i$

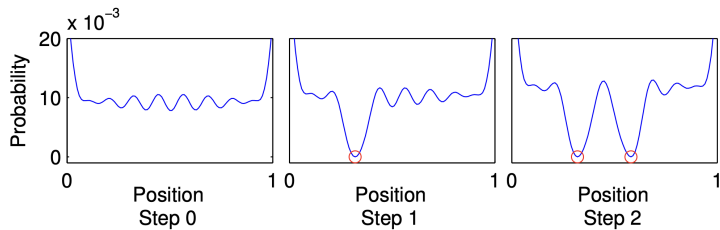
$V \leftarrow V_\perp$ , an orthonormal basis for the subspace of  $V$  orthogonal to  $\mathbf{e}_i$

**end while**

**Output:**  $Y$

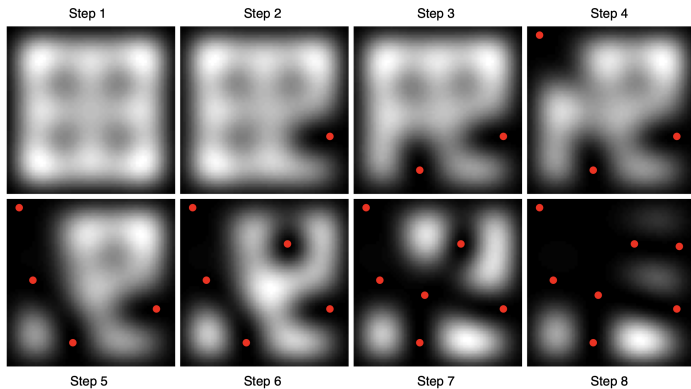
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# Sampling DPPs



(a) Sampling points on an interval

# Sampling DPPs



(b) Sampling points in the plane

# Sampling DPPs

- ▶ Finding the eigendecomposition of  $L$  is  $O(n^3)$ .
- ▶ Sampling algorithm is  $O(n|V|^3)$  for  $V$  the set of eigenvectors selected in phase 1 and each repeated Gram-Schmidt to compute  $V_{\perp}$  is  $O(n|V|^2)$ .



# Dual representation

- ▶ Let  $B$  be the  $D \times N$  matrix with columns  $B_i = \mathbf{q}_i \phi_i$  such that  $L = B^T B$ . Consider the  $D \times D$  matrix

$$C = BB^T.$$

- ▶ Here,  $D$  is the dimension of the diversity feature function  $\phi$ .
- ▶  $D$  is often fixed by design, whereas  $N$  may grow as more items are modeled.

# Dual representation

## Proposition 1

*The non-zero eigenvalues of  $L$  and  $C$  are identical, and the corresponding eigenvectors are related by the matrix  $B$ . That is,*

$$C = \sum_{d=1}^D \lambda_d \hat{v}_d \hat{v}_d^T$$

*is an eigendecomposition of  $C$  if and only if*

$$L = \sum_{d=1}^D \lambda_d \left( \frac{1}{\sqrt{\lambda_d}} B^T \hat{v}_d \right) \left( \frac{1}{\sqrt{\lambda_d}} B^T \hat{v}_d \right)^T$$

*is an eigendecomposition of  $L$ .*

# Dual representation

Proof.

First, assume  $\{\lambda_d, \hat{v}_d\}_{d=1}^D$  is an eigendecomposition of  $C$ . Then,

$$\sum_{d=1}^D \lambda_d \left( \frac{1}{\sqrt{\lambda_d}} B^T \hat{v}_d \right) \left( \frac{1}{\sqrt{\lambda_d}} B^T \hat{v}_d \right)^T = B^T \left( \sum_{d=1}^D \hat{v}_d \hat{v}_d^T \right) B = B^T B = L.$$

Furthermore, we have

$$\begin{aligned} \left\| \frac{1}{\sqrt{\lambda_d}} B^T \hat{v}_d \right\|^2 &= \frac{1}{\lambda_d} (B^T \hat{v}_d)^T (B^T \hat{v}_d) = \frac{1}{\lambda_d} \hat{v}_d^T C \hat{v}_d \\ &= \frac{1}{\lambda_d} \lambda_d \hat{v}_d^T \hat{v}_d = 1, \end{aligned}$$

and

$$\begin{aligned} \left( \frac{1}{\sqrt{\lambda_d}} B^T \hat{v}_d \right)^T \left( \frac{1}{\sqrt{\lambda_{d'}}} B^T \hat{v}_{d'} \right) &= \frac{1}{\sqrt{\lambda_d \lambda_{d'}}} \hat{v}_d^T C \hat{v}_{d'} \\ &= \frac{\sqrt{\lambda_{d'}}}{\sqrt{\lambda_d}} \hat{v}_d^T \hat{v}_{d'} = 0. \end{aligned}$$

A similar argument holds in the other direction when one accounts for the fact  $L = B^T B$  and has rank at most  $D$ . □

## Dual representation and computing

- Normalization: the normalization constant is

$$\det(L + I) = \prod_{d=1}^D (\lambda_d + 1) = \det(C + I),$$

which only takes  $O(D^3)$  time.

- Marginalization: get entries of  $K$  using  $C$ . First get the eigendecomposition  $C = \sum_{d=1}^D \lambda_d \hat{v}_d \hat{v}_d^T$ . Then

$$K_{ij} = \sum_{d=1}^D \frac{\lambda_d}{\lambda_d + 1} \left( \frac{1}{\sqrt{\lambda_d}} B_i^T \hat{v}_d \right)^T \left( \frac{1}{\sqrt{\lambda_d}} B_j^T \hat{v}_d \right).$$

One may therefore obtain the marginal probability of an event in time  $O(D^2)$ . For a  $k$  event, this becomes  $O(D^2 k^2 + k^3)$ . This beats the usual  $O(n^3)$  to translate from  $L$  to  $K$ .

## Dual representation and computing

In general, one may represent the orthonormal set  $V$  in  $\mathbb{R}^n$  using the set  $\hat{V}$  in  $\mathbb{R}^D$  with the mapping

$$V = \{B^T \hat{v} \mid \hat{v} \in \hat{V}\}.$$

One may implicitly obtain linear combinations of vectors in  $V$  by performing actions on their preimages:  $v_1 + v_2 = B^T(\hat{v}_1 + \hat{v}_2)$ .

Moreover,

$$v_1^T v_2 = (B^T \hat{v}_1)^T (B^T \hat{v}_2) = \hat{v}_1^T C \hat{v}_2,$$

so we can compute dot products of elements in  $V$  in time  $O(D^2)$ .

We can implicitly normalize the elements of  $V$  by updating

$$\hat{v} \leftarrow \frac{\hat{v}}{\hat{v}^T C \hat{v}}.$$

# Sampling DPPs

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**Algorithm 1** Sampling from a DPP

---

**Input:** eigendecomposition  $\{(\mathbf{v}_n, \lambda_n)\}_{n=1}^N$  of  $L$

$J \leftarrow \emptyset$

**for**  $n = 1, 2, \dots, N$  **do**

$J \leftarrow J \cup \{n\}$  with prob.  $\frac{\lambda_n}{\lambda_n + 1}$

**end for**

$V \leftarrow \{\mathbf{v}_n\}_{n \in J}$

$Y \leftarrow \emptyset$

**while**  $|V| > 0$  **do**

    Select  $i$  from  $\mathcal{Y}$  with  $\Pr(i) = \frac{1}{|V|} \sum_{\mathbf{v} \in V} (\mathbf{v}^\top \mathbf{e}_i)^2$

$Y \leftarrow Y \cup i$

$V \leftarrow V_\perp$ , an orthonormal basis for the subspace of  $V$  orthogonal to  $\mathbf{e}_i$

**end while**

**Output:**  $Y$

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Can we use the dual representation to speed up the sampling of  $i$  and Gram-Schmidt steps?

## Dual representation and computing

The sampling step is handled thus:

$$\begin{aligned}\Pr(i) &= \frac{1}{|V|} \sum_{v \in V} (v^T e_i)^2 = \frac{1}{|\hat{V}|} \sum_{\hat{v} \in \hat{V}} ((B^T \hat{v})^T e_i)^2 \\ &= \frac{1}{|\hat{V}|} \sum_{\hat{v} \in \hat{V}} (B_i^T \hat{v})^2\end{aligned}$$

The entire distribution may be computed in time  $O(nD|\hat{V}|)$  instead of  $O(n^3)$ .

# Sampling DPPs

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**Algorithm 3** Sampling from a DPP (dual representation)

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**Input:** eigendecomposition  $\{(\hat{\mathbf{v}}_n, \lambda_n)\}_{n=1}^N$  of  $C$

$J \leftarrow \emptyset$

**for**  $n = 1, 2, \dots, N$  **do**

$J \leftarrow J \cup \{n\}$  with prob.  $\frac{\lambda_n}{\lambda_n + 1}$

**end for**

$\hat{V} \leftarrow \left\{ \frac{\hat{\mathbf{v}}_n}{\sqrt{\hat{\mathbf{v}}_n^\top C \hat{\mathbf{v}}_n}} \right\}_{n \in J}$

$Y \leftarrow \emptyset$

**while**  $|\hat{V}| > 0$  **do**

Select  $i$  from  $\mathcal{Y}$  with  $\Pr(i) = \frac{1}{|\hat{V}|} \sum_{\hat{\mathbf{v}} \in \hat{V}} (\hat{\mathbf{v}}^\top B_i)^2$

$Y \leftarrow Y \cup i$

Let  $\hat{\mathbf{v}}_0$  be a vector in  $\hat{V}$  with  $B_i^\top \hat{\mathbf{v}}_0 \neq 0$

Update  $\hat{V} \leftarrow \left\{ \hat{\mathbf{v}} - \frac{\hat{\mathbf{v}}^\top B_i}{\hat{\mathbf{v}}_0^\top B_i} \hat{\mathbf{v}}_0 \mid \hat{\mathbf{v}} \in \hat{V} - \{\hat{\mathbf{v}}_0\} \right\}$

Orthonormalize  $\hat{V}$  with respect to the dot product  $\langle \hat{\mathbf{v}}_1, \hat{\mathbf{v}}_2 \rangle = \hat{\mathbf{v}}_1^\top C \hat{\mathbf{v}}_2$

**end while**

**Output:**  $Y$

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## Quality-diversity representation

In addition to the Gram matrix representation  $L = B^T B$ , we can factor each column  $B_i$  as the product of a 'quality' term  $q_i > 0$  and a normalized 'diversity feature'  $\phi_i \in \mathbb{R}^D$ . Thus,

$$L_{ij} = q_i \phi_i^T \phi_j q_j.$$

If  $q_i$  communicates the 'goodness' of item  $i$ , then

$$S_{ij} = \frac{L_{ij}}{\sqrt{L_{ii} L_{jj}}}.$$

This representation allows one to independently model quality and diversity using the model

$$\Pr_L(Y) \propto \left( \prod_{i \in Y} q_i^2 \right) \det(S_Y)$$

# Conditional DPPs

- ▶ A conditional DPP takes the form of an L-ensemble

$$\Pr_L(Y|X) \propto \det(L_Y(X)).$$

- ▶  $L$  is a positive semi-definite kernel matrix.
- ▶ The normalizing constant takes the form  $\det(L(X) + I)$ .
- ▶ Using the quality-diversity decomposition, we have

$$L_{ij}(X) = q_i(X)\phi_i(X)^T \phi_j(X)q_j(X)$$

for  $q_i > 0$ ,  $\phi_i \in \mathbb{R}^D$  and  $\|\phi_i\| = 1$ .

# Supervised learning

We observe  $\{Y_t, X_t\}_{t=1}^T$  and assume individual  $Y_t$ s generated independently with probabilities

$$\Pr(Y|X, \theta) = \frac{\det(L_Y(X, \theta))}{\det(L(X, \theta) + I)}.$$

Then the log-likelihood takes the form

$$\begin{aligned} \ell(\theta) &= \log \left( \prod_{t=1}^T \Pr(Y_t|X_t, \theta) \right) \\ &= \sum_{t=1}^T \left( \log \det(L_{Y_t}(X_t, \theta)) - \log \det(L(X_t, \theta) + I) \right). \end{aligned}$$

# Supervised learning

Suppose one keeps the feature functions  $\phi_i(X)$  fixed but models the quality scores with the log-linear model

$$q_i(X, \theta) = e^{f_i(X)^T \theta}.$$

Then the probability of a single sample can be written

$$\Pr(Y|X, \theta) = \frac{\det S_Y \prod_{i \in Y} e^{f_i(X)^T \theta}}{\sum_{Y' \subseteq \mathcal{Y}} \det S_{Y'} \prod_{i \in Y'} e^{f_i(X)^T \theta}}.$$

The resulting log-likelihood is convex in  $\theta$ :

$$\ell(\theta) \propto \theta^T \sum_{i \in Y} f_i(X) - \log \sum_{Y' \subseteq \mathcal{Y}} \exp \left( \theta^T \sum_{i \in Y'} f_i(X) \right) \det S_{Y'}(X).$$

## k-DPPs

- ▶ A k-DPP on a discrete set  $\mathcal{Y} = \{1, 2, \dots, N\}$  is a distribution over all sets  $Y \subseteq \mathcal{Y}$  with cardinality  $k$ .
- ▶ A k-DPP is obtained by conditioning a standard DPP on the event that the set  $Y$  has cardinality  $k$ .
- ▶ The k-DPP  $N_L^k$  has probabilities

$$\Pr_L^k(Y) = \frac{\det(L_Y)}{\sum_{|Y'|=k} \det(L_{Y'})}.$$

## k-DPPs: normalization

Define the  $k$ th elementary symmetric polynomial on  $\lambda_1, \dots, \lambda_N$

$$e_k(\lambda_1, \dots, \lambda_N) = \sum_{\substack{J \subseteq \{1, \dots, N\} \\ |J|=k}} \prod_{n \in J} \lambda_n.$$

For example,

$$e_1(\lambda_1, \lambda_2, \lambda_3) = \lambda_1 + \lambda_2 + \lambda_3$$

$$e_2(\lambda_1, \lambda_2, \lambda_3) = \lambda_1 \lambda_2 + \lambda_1 \lambda_3 + \lambda_2 \lambda_3$$

$$e_3(\lambda_1, \lambda_2, \lambda_3) = \lambda_1 \lambda_2 \lambda_3.$$

### Proposition 2

*The normalizing constant for a  $k$ -DPP is*

$$Z_k = \sum_{|Y'|=k} \det(L_{Y'}) = e_k(\lambda_1, \dots, \lambda_N),$$

*where  $\lambda_n$  are the eigenvalues of  $L$ .*

## k-DPPs: normalization

Proof.

Recalling that

$$\sum_{Y \subseteq \mathcal{Y}} \det(L_Y) = \det(L + I),$$

we know

$$\sum_{|Y'|=k} \det(L_{Y'}) = \det(L + I) \sum_{|Y'|=k} \Pr_L(Y').$$

Then, because every DPP is a mixture of elementary DPPs:

$$\begin{aligned} \det(L + I) \sum_{|Y'|=k} \Pr_L(Y') &= \frac{\det(L + I)}{\det(L + I)} \sum_{|Y'|=k} \sum_{J \subseteq \{1, \dots, N\}} \Pr^{V_J}(Y') \prod_{n \in J} \lambda_n \\ &= \sum_{|J|=k} \sum_{|Y'|=k} \Pr^{V_J}(Y') \prod_{n \in J} \lambda_n \\ &= \sum_{|J|=k} \prod_{n \in J} \lambda_n. \end{aligned}$$

□

# Computing elementary symmetric polynomials

Use the shorthand  $e_k^N = e_k(\lambda_1, \dots, \lambda_N)$ , we have the recursion

$$e_k^N = e_k^{N-1} \lambda_N e_{k-1}^{N-1}.$$

Thus, the following algorithm computes  $e_k^N$  in time  $O(Nk)$ .

---

**Algorithm 7** Computing the elementary symmetric polynomials

---

**Input:**  $k$ , eigenvalues  $\lambda_1, \lambda_2, \dots, \lambda_N$

$e_0^n \leftarrow 1 \quad \forall n \in \{0, 1, 2, \dots, N\}$

$e_l^0 \leftarrow 0 \quad \forall l \in \{1, 2, \dots, k\}$

**for**  $l = 1, 2, \dots, k$  **do**

**for**  $n = 1, 2, \dots, N$  **do**

$e_l^n \leftarrow e_l^{n-1} + \lambda_n e_{l-1}^{n-1}$

**end for**

**end for**

**Output:**  $e_k(\lambda_1, \lambda_2, \dots, \lambda_N) = e_k^N$

---



## k-DPPs: sampling

- ▶ One may use a (slow) rejection sampling approach, sampling DPPs and discarding those for which  $|Y| \neq k$ .
- ▶ It is more efficient to first recognize that, when  $|Y| = k$

$$\Pr_L^k(Y) = \frac{\det(L + I)}{e_k^N} \Pr_L(Y)$$

and therefore

$$\Pr_L^k(Y) = \frac{1}{e_k^N} \sum_{|J|=k} \Pr^{V_J}(Y) \prod_{n \in J} \lambda_n.$$

- ▶ A k-DPP is also a mixture of elementary DPPs! So *if* we can sample  $k$  eigenvalues, we can then use the mixture of elementary DPPs to generate samples.

## k-DPPs: sampling

The following  $O(Nk)$  algorithm samples sets of  $k$  eigenvalues according to desired probabilities

$$\Pr(J) = \frac{\mathbf{1}\{|J| = k\}}{e_k^N} \prod_{n \in J} \lambda_n.$$

---

**Algorithm 8** Sampling  $k$  eigenvectors

---

**Input:**  $k$ , eigenvalues  $\lambda_1, \lambda_2, \dots, \lambda_N$

compute  $e_l^n$  for  $l = 0, 1, \dots, k$  and  $n = 0, 1, \dots, N$  (Algorithm 7)

$J \leftarrow \emptyset$

$l \leftarrow k$

**for**  $n = N, \dots, 2, 1$  **do**

**if**  $l = 0$  **then**

**break**

**end if**

**if**  $u \sim U[0, 1] < \lambda_n \frac{e_{l-1}^{n-1}}{e_l^n}$  **then**

$J \leftarrow J \cup \{n\}$

$l \leftarrow l - 1$

**end if**

**end for**

**Output:**  $J$

---

## k-DPPs: marginalization

Recall that for a general L-ensemble, we have

$$\begin{aligned}\Pr_L(B \subseteq Y | A \subseteq Y) &= \det \left( [I - [(L + I_{A^c})^{-1}]_{A^c}]_B \right) \\ &= \det(L_B^A).\end{aligned}$$

## k-DPPs: marginalization

k-DPPs are not DPPs and do not have a marginal kernel. But for  $|A| \leq k$ , we have:

$$\begin{aligned}\Pr_L^k(A \subseteq Y) &= \sum_{\substack{|Y'|=k-|A| \\ Y' \cap A = \emptyset}} \Pr_L^k(Y' \cup A) \\ &= \frac{\det(L + I)}{Z_k} \sum_{\substack{|Y'|=k-|A| \\ Y' \cap A = \emptyset}} \Pr_L(Y' \cup A) \\ &= \frac{\det(L + I)}{Z_k} \sum_{\substack{|Y'|=k-|A| \\ Y' \cap A = \emptyset}} \Pr_L(Y = Y' \cup A | A \subseteq Y) \Pr_L(A \subseteq Y) \\ &= \frac{Z_{k-|A|}^A \det(L + I)}{Z_k \det(L^A + I)} \Pr_L(A \subseteq Y),\end{aligned}$$

where

$$Z_{k-|A|}^A = \det(L^A + I) \sum_{\substack{|Y'|=k-|A| \\ Y' \cap A = \emptyset}} \Pr_L(Y = Y' \cup A | A \subseteq Y) = \sum_{\substack{|Y'|=k-|A| \\ Y' \cap A = \emptyset}} \det(L_{Y'}^A)$$

is the normalizing constant for the  $(k - |A|)$ -DPP with kernel  $L^A$ .

## k-DPPs: marginalization

Thus, the marginal probabilities for a k-DPP are the same as those of the DPP with the same kernel but properly renormalized. By observing that

$$\frac{\det(L^A)}{\det(L + I)} = \frac{\Pr_L(A \subseteq Y)}{\det(L^A + I)},$$

(since  $1/\det(L^A + I)$  is the probability of observing nothing else conditioned on  $A$ ), the equation simplifies further:

$$\begin{aligned}\Pr_L^k(A \subseteq Y) &= \frac{Z_{k-|A|}^A}{Z_k} \frac{\det(L + I)}{\det(L^A + I)} \Pr_L(A \subseteq Y) \\ &= \frac{Z_{k-|A|}^A}{Z_k} \det(L^A) = Z_{k-|A|}^A \Pr_L^k(A).\end{aligned}$$

Computing such a probability is  $O((N - |A|)^3)$  and very inefficient for  $|A|$  small.

## k-DPPs: singleton marginals

First, write the marginal probability of an item  $i$  using elementary DPPs:

$$\Pr_L^k(i \in Y) = \frac{1}{e_k^N} \sum_{|J|=k} \Pr^{V_J}(i \in Y) \prod_{n' \in J} \lambda_{n'}$$

But the marginal kernel of an elementary DPP is  $\sum_{n \in J} v_n v_n^T$ , so this becomes:

$$\begin{aligned} \Pr_L^k(i \in Y) &= \frac{1}{e_k^N} \sum_{|J|=k} \left( \sum_{n \in J} (e_i^T v_n)^2 \right) \prod_{n' \in J} \lambda_{n'} \\ &= \frac{1}{e_k^N} \sum_{n=1}^N (e_i^T v_n)^2 \sum_{\substack{J \supset \{n\} \\ |J|=k}} \prod_{n' \in J} \lambda_{n'} \\ &= \sum_{n=1}^N (e_i^T v_n)^2 \lambda_n \frac{e_{k-1}^{-n}}{e_k^N} \end{aligned}$$

If we have the eigendecomposition of  $L$  and know the values  $e_{k-1}^{-n}/e_k^N$ , then we can obtain all singleton marginals in time  $O(N^2)$ .  $e_k^N$  can be computed in time  $O(Nk)$  and all  $e_{k-1}^{-n}$  can be computed in time  $O(N^2k)$ . This can be improved to  $O(N \log(N)k)$ .

## k-DPPs: conditioning

For  $|A| + |B| = k$ ,

$$\begin{aligned}\Pr_L^k(Y = A \cup B | A \subseteq Y) &\propto \Pr_L^k(Y = A \cup B) \\ &\propto \Pr_L(Y = A \cup B) \\ &\propto \Pr_L(Y = A \cup B | A \subseteq Y) \\ &\propto \det(L_B^A).\end{aligned}$$

So the conditional k-DPP is a  $(k - |A|)$ -DPP.